## **ROS 2 on Autonomous Vehicles**

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### Introduction

Autonomous vehicles will... ... give hours back to commuters, ...change the way the world is connected, ...disrupt industries, and ...generate lots of value.

Autonomous vehicles are big robots with...

- ...sensors,
- ...actuation, and
- ...lots of algorithms,
- ...and they can cause a lot of damage.



### Can we make an autonomous vehicle<sup>1</sup> using ROS 2?

Yes<sup>2</sup>

<sup>1)</sup> a large robotic system in a safety critical application. <sup>2)</sup> with caveats.









#### Round About



### Our Autonomous Driving Setup







### Our ROS 2 - ROS 1 Setup





### Recap: ROS 2 vs. ROS 1

ROS 2 has all the core features needed to build a large robot

- Node API, topics, services
- Parameter server
- Command line introspection
- Composition
- **TF**

#### In addition to extra features missing from ROS 1

- Deterministic roslaunch
- Rclcpp\_lifecycle
- DDS (best effort and reliable QoS)
- Data Security
- Layered architecture ullet

#### ros1\_bridge

#### Why we still need ROS 1

- Legacy Algorithms
- Rviz
- Rosbag
- Rqt\_graph
- Rqt\_plot
- Gazebo  $\bullet$
- Console Logging (to file, rosout topic)



### Limitations of ROS 2 for Autonomous Driving

ROS 2 is missing features needed for safety-critical applications

- 1. Hard Real-Time
  - OS Primitives (memory, synchronization)
  - Real Time Logging
  - Waitsets
  - Large Memory Support
- 2. Robustness and Security
  - Managed System
  - OS Security
- 3. Testing and Certification

Apex.OS is an automotive ROS 2 for safety-critical applications



### Automotive ROS 2: Hard Real Time

#### Deterministic resource usage and runtime is necessary for a safety critical system

- Memory
- Threads
- Blocking calls

#### ROS 2 is still too dynamic for hard real-time

#### 1. Memory

- Allocation on subscription
- std::string
- std::vector
- std::exception
- 2. Blocking calls
  - fprintf
  - fwrite
- 3. Non-RT DDS Implementation

To bridge the gap to hard real-time

- No resource allocation during runtime ullet
- All operations are finite and bounded
- All potentially blocking calls have timeouts



## Automotive ROS 2: Real Time Logging

Printing console is a nondeterministic blocking call

// Trace of RCLCPP\_INFO call: #define APEX\_PRINT(...) \ RCLCPP\_INFO(logger, "foo"); apex::console::print( \ RCUTILS\_LOG\_INFO\_NAMED(logger.get\_name(), "foo"); static\_cast<uint32\_t>(\_\_LINE\_\_), \ RCUTILS\_LOG\_COND\_NAMED(...);  $FILE_{}$ // ... VA ARGS rcutils\_log(...); APEX\_PRINT("Debug float value", 32.23F); (\*output\_handler) (location, severity, name ? name : "", now, format, &args); output\_handler = g\_rcutils\_logging\_output\_handler; 149411 apex\_console\_logging | 2018-08-30 17:15:54.3195151 Version: 0.0.0 g\_rcutils\_logging\_output\_handler = 14942l apex\_console\_logging | 2018-08-30 17:15:54.319520l Formal build: No rcutils\_logging\_console\_output\_handler(...); 14943l apex\_console\_logging | 2018-08-30 17:15:54.319521l Debug float value: :+32.23 14944l apex\_console\_logging | 2018-08-30 17:15:55.319628l Debug integer value: :-32 // Calls 14945l apex\_console\_logging | 2018-08-30 17:15:56.319741l This is a debug message fprintf(...);

A purpose-built real-time logger was built instead

• Logging call uses deterministic atomic

operations

• Writes to a self-healing, fail-resistent ring-

buffer in shared memory

• Buffer can be flushed with minimal

overhead



#### Automotive ROS 2: Waitsets

Waitsets better lend themselves to a deterministic Callbacks are the primary mechanism by which ROS handles the receipt of execution order and error handling interprocess communication

**const auto** subscriber1\_ptr = node\_ptr->create\_subscriber<std\_msgs::msg::String>( "Topic1", bar); **const auto** subscriber2\_ptr = node\_ptr->create\_subscriber<geometry\_msgs::msg::PointStamped>( "Topic2", foo);

// foo and bar get executed in an arbitrary order rclcpp::spin(node\_ptr);

```
rclcpp::Node node("Node");
auto sub1 =
 node.create_subscriber<std_msgs::msg::String>(
 "Topic1");
auto sub2 =
 node.create_subscriber<geometry_msgs::msg::PointStamped>(
 "aTopic2");
rclcpp::Waitset<2> ws({sub1, sub});
```

```
// Wait for 5 seconds.
ws.wait(5s);
```

```
auto msgs1 = sub1->take();
if(msgs1) {
 // always update to latest msgs2 if available
 // before acting on msg1
 auto msgs2 = sub2->take();
 if(msgs2) {
  handle_sample(msg2.data());
```

```
handle_sample(msg1.data());
} else {
// react to not receiving msgs1 in time
```



### Automotive ROS 2: Large Memory Support

#### The maximum size of a UDP packet is 64kB

- Messages larger than 64kB require fragmentation
- Large messages are slower to transmit •
- Exchanging pointers (8 B) to memory locations in shared memory is significantly faster for large data



/\* SHM Publish \*/ // Initialize apex::shared\_memory::ShmArray<BigMsg> shm\_pub(num\_frames, topic.c\_str()); **const auto** pub\_ptr = node\_ptr->create\_publisher<std\_msgs::msg::Uint64>(topic); // publish: write message to shared memory

BigMsg big\_msg; **const** uint64\_t frame\_num = **0**U; shm\_pub[frame\_num] = big\_msg; // publish: send frame number via DDS std\_msgs::msg::Uint64 msg; msg.data = frame\_num; pub\_ptr->publish(msg);

/\* SHM Subscribe \*/ // Initialize const apex::shared\_memory::ShmArray<BigMsg> shm\_sub(num\_frames, topic.c\_str());

auto cb = [&](const std\_msgs::msg::Uint64::SharedPtr msg) { // copy large message to local context // could also manipulate in shared memory for zero copy local\_big\_msg = shm\_sub[msg->data];

**const auto** sub\_ptr = node\_ptr->create\_publisher<std\_msgs::msg::Uint64>(topic, cb);



### Automotive ROS 2: Managed System

Deterministic startup order of nodes is important for large systems

- ROS 2 launch (Python) provides this capability
- ROS 2's managed nodes allow individual nodes to react to failures

ROS 2 lacks mechanisms for the whole system to react to node failures

- Heartbeat (detect silent failures)
- Lifecycle Manager (coordinate system level responses)
- Shadow nodes (instant failure response for critical systems)
  - Consensus

ullet



### Automotive ROS 2: Security

ROS 2 exposes three kinds of security from DDS

- Message encryption
- Authentication
- Access Control
- Data Tagging
- Logging

This is insufficient to guard against corrupted or malicious binaries

- Memory hoggers
- CPU stressors
- Tailor-made DDS participants

ROS 2 also lacks key mechanisms such as • Secure over-the-air (OTA) updates • Secure key storage Integration with existing security infrastructure



#### Automotive ROS 2: Testing and Certification How do you prove code is safe?

#### Follow a functional safety standard (ISO 26262):

- 1. Analyze your use case
- 2. Follow a process

  - Follow a coding standard
  - Analyze your code
  - Do code reviews
- 3. Write tests
  - Unit, integration, full stack, stress, fault, injection, requirements
  - SIL, HIL (every supported ECU, sensor)
  - Line, branch, MC/DC coverage







### Giving Back

Tools	Static Exceptions
	Agile Development Environr
	(Inter-Process Communicati
	ament_pclint
	AutowareAuto
ROS 2 Features	YAML Parameter Parser
Bugs	osrf_testing_tools_cpp
	orocos_kinematics_dynamic
	rcutils

# nment tion) Performance Test

ics



#### Building Algorithms for Safety Critical Applications Case Study: ROS 1 Velodyne Driver

As is:	Changes:
Memory allocation per packet	All memory
Nodelets as a proxy for threading	Threading w
OS-specific system calls	OS-agnostic
No failure handling	Notifies use
Coarse point cloud discretization	Firing-level
Non-RT logging	RT logging
Integration tests	Unit, integra



- allocation on startup
- with controllable stack size, priority, etc.
- c system calls
- er on UDP timeout
- point cloud discretization

ation, stress tests



### Conclusion

#### ROS 1 is great

- Lots of tools and algorithms
- Community
- But can never be automotive grade ROS 2 is even greater
- Enough features for serious development
- API is stable enough
- Can work with ROS 1
- Can be automotive grade

#### **Apex.OS** is automotive grade ROS 2

- Real-time
- Secure
- ISO 26262 certified

## Apex.A<sup>®</sup> is hiring!

- We are actively recruiting developers and engineers:
- Framework
- Embedded
- Security
- Certification
- Algorithms
- If you are interested to learn more, talk to us
- at ROSCon or apply at www.apex.ai



## ROS 1 PubSub vs. Apex.OS PubSub







Rate & Latency between nodes under stress (Run-2)	
400	<ul> <li>Pub_Sub_Lat(ms)</li> <li>Rate (ms)</li> </ul>
300	
200	
100 ·····///////////////////////////////	

